

## Computer Science 456/656 Spring 2013

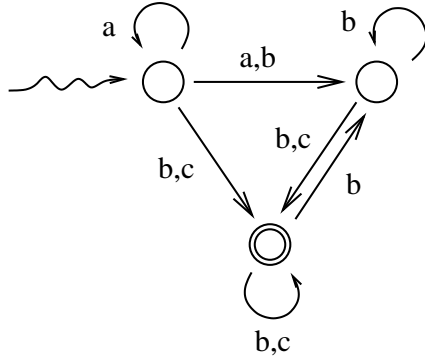
Practice for the First Examination, February 28, 2013

**The entire practice examination is 455 points. The real exam will be much shorter.**

1. True or False. [5 points each]

- (a) ..... Every subset of a regular language is regular.
- (b) ..... Let  $L$  be the language over  $\Sigma = \{a, b\}$  consisting of all strings of the form  $a^m b^n$ , where  $m, n \geq 0$ . Then  $L$  is a regular language.
- (c) ..... The complement of every regular language is regular.
- (d) ..... The Kleene closure of every context-free language is context-free.
- (e) ..... If a language has an unambiguous context-free grammar, then it is accepted by some deterministic push-down automaton.
- (f) ..... If a language has an ambiguous context-free grammar, then it is not accepted by any deterministic push-down automaton.
- (g) ..... There is a PDA that accepts all valid C++ programs.
- (h) ..... The intersection of any two regular languages is regular.
- (i) ..... The language consisting of all base 7 numerals for positive integers  $n$  such that  $n \% 3 = 2$  is regular.
- (j) ..... The intersection of any two context-free languages is context-free.
- (k) ..... Let  $L$  be the language over  $\Sigma = \{a, b\}$  consisting of all strings of the form  $a^m b^n c^m$ , where  $m, n \geq 0$ . Then  $L$  is a context-free language.
- (l) ..... Let  $L$  be the language over  $\Sigma = \{a, b\}$  consisting of all strings of the form  $a^m b^n$ , where  $m \geq n$ . Then  $L$  is a context-free language.
- (m) ..... The complement of every context-free language is context-free.
- (n) ..... The union of any two context-free languages is context-free.
- (o) ..... If a language has a context-free grammar, then it is accepted by some push-down automaton.
- (p) ..... Every context-free language has an unambiguous context-free grammar.
- (q) ..... Every language that has an unambiguous context-free grammar is accepted by some DPDA.
- (r) ..... The intersection of any two context-free languages is context-free.
- (s) ..... Every deterministic machine is a non-deterministic machine.
- (t) ..... The language consisting of all base 2 numerals for integer powers of 2 is regular.
- (u) ..... There is a DPDA that accepts the language of all palindromes over the binary alphabet  $\{0, 1\}$ .

2. [25 points] Draw an NFA with five states which accepts the language described by the regular expression  $(0 + 1)^*0(0 + 1)(0 + 1)(0 + 1)$
3. [25 points] Write a regular expression for the language accepted by the following NFA. If your answer is unnecessarily long by a wide margin, I might mark it wrong even if it's right.



Find a Regular Expression

4. [20 points] Let  $G$  be the context-free grammar given below.
  - (a)  $S \rightarrow a$
  - (b)  $S \rightarrow wS$
  - (c)  $S \rightarrow iS$
  - (d)  $S \rightarrow iSeS$

Prove that  $G$  is ambiguous by writing two different **leftmost** derivations for the string  $iwiAEA$ . [If you simply show two different parse trees, you are not following instructions.]

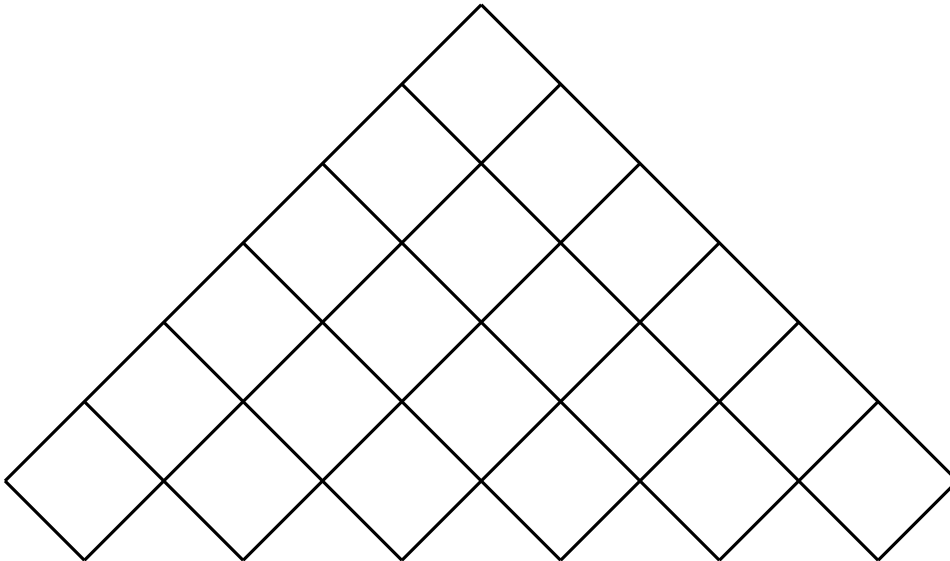
5. [30 points] Design a PDA that accepts the language  $L = \{a^nbc^n : n \geq 0\}$ .
6. [30 points] Give a context-free grammar for the language of all strings over  $\{0, 1\}$  of the form  $0^m1^n$  where  $n \neq m$ .
7. [30 points] The following context-free grammar  $G$  is ambiguous. Give an equivalent unambiguous grammar.

- The terminal alphabet of  $G$  is  $\{a, b, c, (, ), +, -, *\}$ .
- $G$  has only one variable, namely the start symbol  $E$ .
- The productions of  $G$  are as follows:
  - (a)  $E \rightarrow E + E$
  - (b)  $E \rightarrow E - E$
  - (c)  $E \rightarrow E * E$
  - (d)  $E \rightarrow (E)$
  - (e)  $E \rightarrow a$
  - (f)  $E \rightarrow b$
  - (g)  $E \rightarrow c$

8. [30 points] Let  $L$  be the language generated by the Chomsky Normal Form (CNF) grammar given below.

- (a)  $S \rightarrow a$
- (b)  $E \rightarrow a$
- (c)  $S \rightarrow LA$
- (d)  $E \rightarrow LA$
- (e)  $L \rightarrow ($
- (f)  $A \rightarrow ER$
- (g)  $R \rightarrow )$
- (h)  $S \rightarrow PE$
- (i)  $E \rightarrow PE$
- (j)  $S \rightarrow EE$
- (k)  $E \rightarrow EE$
- (l)  $P \rightarrow EQ$
- (m)  $Q \rightarrow +$

Use the CYK algorithm to prove that the string  $a(a + a)$  is a member of  $L$ . Use the figure below for your work.



9. [15 points] State the pumping lemma for regular languages.

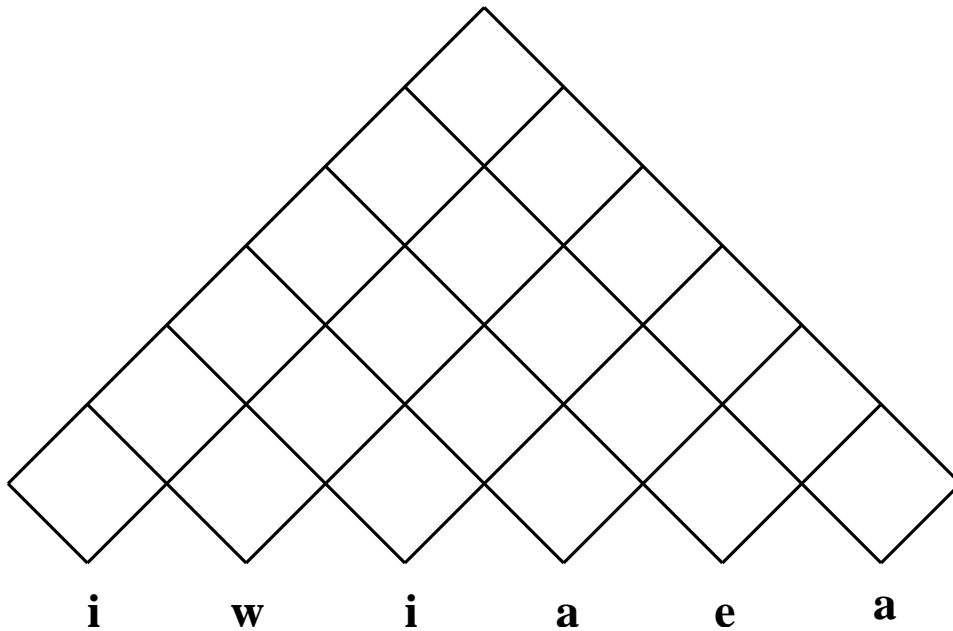
10. [15 points] State the pumping lemma for context-free languages.

11. Let  $L$  be the language generated by the context-free grammar given in the first box. The second box contains a Chomsky Normal Form (CNF) grammar that also generates  $L$ .

- (a)  $S \rightarrow a$
- (b)  $S \rightarrow wS$
- (c)  $S \rightarrow iS$
- (d)  $S \rightarrow iSeS$

- (a)  $S \rightarrow WS$
- (b)  $S \rightarrow IS$
- (c)  $S \rightarrow AB$
- (d)  $A \rightarrow IS$
- (e)  $B \rightarrow ES$
- (f)  $S \rightarrow a$
- (g)  $I \rightarrow i$
- (h)  $W \rightarrow w$
- (i)  $E \rightarrow e$

- (a) [20 points] Use the CYK algorithm to prove that the string  $iwiaea$  is a member of  $L$ . Use the figure below for your work.
- (b) [20 points] By looking at your work carefully, you can determine that the CNF grammar given above is ambiguous. Write two different parse trees for  $iwiaea$ , using the CNF grammar.



12. [30 points] Consider the NFA whose transition diagram is drawn below, where the input alphabet is  $\{a, b, c\}$ . Draw the transition diagram of an equivalent minimal DFA. Show your steps.

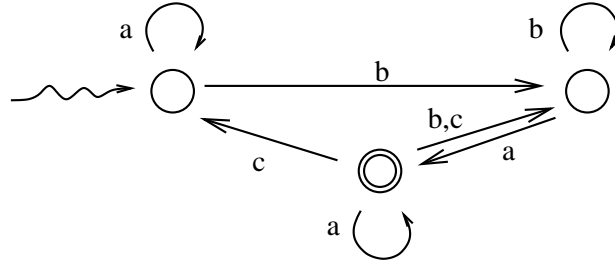


Figure 1: Find a minimal DFA equivalent to this NFA

13. [30 points] Let  $L = \{w \in \{a, b\}^* \mid \#_a(w) = 2\#_b(w)\}$ , here  $\#_a(w)$  denotes the number of instances of the symbol  $a$  in the string  $w$ . For example,  $aaababaaabba \in L$ , because that string has the twice as many  $a$ 's as  $b$ 's. Give a context-free grammar for  $L$ . Your grammar may be ambiguous.

14. [30 points]

1.  $S \rightarrow \epsilon$

2.  $S \rightarrow a_2S_3b_4S_5$

|   | $a$ | $b$ | eof  | $S$ |
|---|-----|-----|------|-----|
| 0 |     |     |      |     |
| 1 |     |     | halt |     |
| 2 |     |     |      |     |
| 3 |     |     |      |     |
| 4 |     |     |      |     |
| 5 |     |     |      |     |

Complete the ACTION and GOTO tables of an LALR parser for the grammar given above. This grammar unambiguously generates the “balanced parentheses” language, where  $a$  represents a left parenthesis, and  $b$  represents a right parenthesis. Example strings include  $\epsilon$ ,  $ab$ ,  $aabb$ ,  $abab$ , and  $aabbab$ .