Administrative Notes

- Note: New homework instructions starting with HW03
- Homework is due at the beginning of class
- Homework must be organized, legible (messy is not), and stapled to be graded



Some Definitions

- Complement: variable with a bar over it $\overline{A}, \overline{B}, \overline{C}$
- Literal: variable or its complement
 A, A, B, B, C, C
- Implicant: product of literals
 ABC, AC, BC
- Minterm: product that includes all input variables

$AB\overline{C}, A\overline{B}\overline{C}, ABC$

Maxterm: sum that includes all input variables
 (A+B+C), (A+B+C), (A+B+C)



Canonical Sum-of-Products (SOP) Form

- All equations can be written in SOP form
- Each row has a **minterm**
- A minterm is a product (AND) of literals
- Each minterm is TRUE for that row (and only that row)

| | | | | minterm |
|---|---|---|---------|---------------|
| A | B | Y | minterm | name |
| 0 | 0 | 0 | A B | m_0 |
| 0 | 1 | 1 | Ā B | m_1° |
| 1 | 0 | 0 | AB | m_2 |
| 1 | 1 | 1 | ΑB | m_{3} |



Canonical Sum-of-Products (SOP) Form

- All equations can be written in SOP form
- Each row has a **minterm**
- A minterm is a product (AND) of literals
- Each minterm is TRUE for that row (and only that row)
- Form function by ORing minterms where the output is TRUE

| | | | | minterm |
|---|---|---|---------|---------------|
| A | B | Y | minterm | name |
| 0 | 0 | 0 | A B | m_0 |
| 0 | 1 | 1 | Ā B | $\tilde{m_1}$ |
| 1 | 0 | 0 | AB | m_2 |
| 1 | 1 | 1 | ΑB | m_3 |

 $Y = \mathbf{F}(A, B) =$





Canonical Sum-of-Products (SOP) Form

- All equations can be written in SOP form
- Each row has a **minterm**
- A minterm is a product (AND) of literals
- Each minterm is TRUE for that row (and only that row)
- Form function by ORing minterms where the output is TRUE
- Thus, a sum (OR) of products (AND terms)

| | | | | minterm |
|---|---|---|---------|---------|
| A | B | Y | minterm | name |
| 0 | 0 | 0 | A B | m_0 |
| 0 | 1 | 1 | Ā B | m_1 |
| 1 | 0 | 0 | A B | m_2 |
| 1 | 1 | 1 | ΑB | m_3 |

 $Y = F(A, B) = AB + AB = \Sigma(m_1, m_3)$



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Chapter 2 <13>

SOP Example

- Steps:
- Find minterms that result in Y=1
- Sum "TRUE" minterms

| Α | В | Y |
|---|---|---|
| 0 | 0 | 1 |
| 0 | 1 | 1 |
| 1 | 0 | 0 |
| 1 | 1 | 0 |

 $Y = \mathbf{F}(A, B) =$



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Chapter 2 <14>

Aside: Precedence

- AND has precedence over OR
- In other words:
 - AND is performed **before** OR

- Example:
 - $Y = \overline{A} \cdot B + A \cdot B$
 - Equivalent to:
 - $Y = (\overline{A}B) + (AB)$



Canonical Product-of-Sums (POS) Form

- All Boolean equations can be written in POS form
- Each row has a **maxterm**
- A maxterm is a sum (OR) of literals
- Each maxterm is FALSE for that row (and only that row)

| | | | | maxterm |
|----|---|---|-------------------------------|----------------|
| Α | B | Y | maxterm | name |
| 0 | 0 | 0 | A + B | M _o |
| 0 | 1 | 1 | $A + \overline{B}$ | M_1 |
| (1 | 0 | 0 | <u>A</u> + B | M_2 |
| 1 | 1 | 1 | $\overline{A} + \overline{B}$ | M_3 |



Canonical Product-of-Sums (POS) Form

- All Boolean equations can be written in POS form
- Each row has a **maxterm**
- A maxterm is a sum (OR) of literals
- Each maxterm is FALSE for that row (and only that row)
- Form function by ANDing the maxterms for which the output is FALSE
- Thus, a product (AND) of sums (OR terms)

| | | | | maxterm |
|------------|---------|-------------|-------------------------------|------------------------------|
| Α | B | Y | maxterm | name |
| 0 | 0 | 0 | A + B | M_0 |
| 0 | 1 | 1 | $A + \overline{B}$ | M_1 |
| (1 | 0 | 0 | <u>A</u> + B | M_2 |
| 1 | 1 | 1 | $\overline{A} + \overline{B}$ | M_3 |
| <i>Y</i> = | $= M_0$ | $\cdot M_2$ | = (A + B) | $) \cdot (\overline{A} + B)$ |



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Chapter 2 <17>

SOP and POS Comparison

- Sum of Products (SOP)
 - Implement the "ones" of the output
 - Sum all "one" terms \rightarrow OR results in "one"
- Product of Sums (POS)
 - Implement the "zeros" of the output
 - Multiply "zero" terms \rightarrow AND results in "zero"



Boolean Equations Example

- You are going to the cafeteria for lunch
 - You will eat lunch (E=1)
 - If it's open (O=1) and
 - If they're not serving corndogs (C=0)
- Write a truth table for determining if you will eat lunch (E).



Boolean Equations Example

- You are going to the cafeteria for lunch
 - You will eat lunch (E=1)
 - If it's open (O=1) and
 - If they're not serving corndogs (C=0)
- Write a truth table for determining if you will eat lunch (E).



• SOP – sum-of-products

| C |) C | E | minterm |
|---|-----|---|-----------------------------|
| 0 | 0 | | $\overline{O} \overline{C}$ |
| 0 | 1 | | O C |
| 1 | 0 | | ΟC |
| 1 | 1 | | ΟC |

• POS – product-of-sums

| 0 | С | Ε | maxterm |
|---|---|---|-------------------------------|
| 0 | 0 | | 0 + C |
| 0 | 1 | | $O + \overline{C}$ |
| 1 | 0 | | <u> </u> |
| 1 | 1 | | $\overline{O} + \overline{C}$ |



• SOP – sum-of-products

| 0 | С | Е | minterm |
|----|---|---|----------------|
| 0 | 0 | 0 | |
| 0 | 1 | 0 | O C |
| (1 | 0 | 1 | \overline{O} |
| 1 | 1 | 0 | ΟC |

• POS – product-of-sums

| 0 | С | Е | maxterm |
|---|---|---|-------------------------------|
| 0 | 0 | 0 | O + C |
| 0 | 1 | 0 | $O + \overline{C}$ |
| 1 | 0 | 1 | <u> </u> |
| 1 | 1 | 0 | $\overline{O} + \overline{C}$ |



• SOP – sum-of-products

| | minterm | Е | С | 0 |
|-----------------|-----------------|---|---|---|
| | | 0 | 0 | 0 |
| E = OC | O C | 0 | 1 | 0 |
| $= \Sigma(m_2)$ | $O\overline{C}$ | 1 | 0 | 1 |
| <u>/</u> / | ΟC | 0 | 1 | 1 |

• POS – product-of-sums

| 0 | С | Е | maxterm |
|----|---|---|-------------------------------|
| 0 | 0 | 0 | O + C |
| 0 | 1 | 0 | $O + \overline{C}$ |
| 1 | 0 | 1 | <u> </u> |
| (1 | 1 | 0 | $\overline{O} + \overline{C}$ |



• SOP – sum-of-products

| 0 | С | Е | minterm | |
|---|---|---|---------|----------------|
| 0 | 0 | 0 | | |
| 0 | 1 | 0 | O C | E = OC |
| 1 | 0 | 1 | 0 C | $=\Sigma(m_2)$ |
| 1 | 1 | 0 | ΟC | -(2) |

POS – product-of-sums

| 0 | С | Е | maxterm |
|---|---|---|-------------------------------|
| 0 | 0 | 0 | O + C |
| 0 | 1 | 0 | $O + \overline{C}$ |
| 1 | 0 | 1 | <u> </u> |
| 1 | 1 | 0 | $\overline{O} + \overline{C}$ |

 $E = (O + C)(O + \overline{C})(\overline{O} + \overline{C})$ $= \Pi(M0, M1, M3)$



Boolean Algebra

- Axioms and theorems to simplify Boolean equations
- Like regular algebra, but simpler: variables have only two values (1 or 0)
- **Duality** in axioms and theorems:
 - ANDs and ORs, 0's and 1's interchanged



Axiom

| A1 | $B = 0$ if $B \neq 1$ |
|----|---------------------------------|
| A2 | $\overline{0} = 1$ |
| A3 | $0 \bullet 0 = 0$ |
| A4 | 1 • 1 = 1 |
| A5 | $0 \bullet 1 = 1 \bullet 0 = 0$ |
| | |



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Chapter 2 <26>

Duality

Duality in Boolean axioms and theorems:

- ANDs and ORs, 0's and 1's interchanged



Axiom

| A1 | $B = 0$ if $B \neq 1$ |
|----|---------------------------------|
| A2 | $\overline{0} = 1$ |
| A3 | $0 \bullet 0 = 0$ |
| A4 | 1 • 1 = 1 |
| A5 | $0 \bullet 1 = 1 \bullet 0 = 0$ |
| | |



| | Axiom | | Dual |
|----|---------------------------------|-----|-----------------------|
| A1 | $B = 0$ if $B \neq 1$ | A1′ | $B = 1$ if $B \neq 0$ |
| A2 | $\overline{0} = 1$ | A2′ | $\overline{1} = 0$ |
| A3 | $0 \bullet 0 = 0$ | A3′ | 1 + 1 = 1 |
| A4 | 1 • 1 = 1 | A4′ | 0 + 0 = 0 |
| A5 | $0 \bullet 1 = 1 \bullet 0 = 0$ | A5′ | 1 + 0 = 0 + 1 = 1 |

Dual: Exchange: • and + 0 and 1



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Chapter 2 <29>

| | Axiom | | Dual | Name |
|----|---------------------------------|-----|-----------------------|--------------|
| A1 | $B = 0$ if $B \neq 1$ | A1′ | $B = 1$ if $B \neq 0$ | Binary field |
| A2 | $\overline{0} = 1$ | A2′ | $\overline{1} = 0$ | NOT |
| A3 | $0 \bullet 0 = 0$ | A3′ | 1 + 1 = 1 | AND/OR |
| A4 | 1 • 1 = 1 | A4′ | 0 + 0 = 0 | AND/OR |
| A5 | $0 \bullet 1 = 1 \bullet 0 = 0$ | A5′ | 1 + 0 = 0 + 1 = 1 | AND/OR |

Dual: Exchange: • and + 0 and 1



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Chapter 2 <30>

Basic Boolean Theorems

| | Theorem |
|----|-------------------------------|
| T1 | $B \bullet 1 = B$ |
| T2 | $B \bullet 0 = 0$ |
| Τ3 | $B \bullet B = B$ |
| T4 | $\overline{\overline{B}} = B$ |
| T5 | $B \bullet \overline{B} = 0$ |



Basic Boolean Theorems: Duals

| | Theorem | | Dual | Name |
|----|------------------------------|-------------------------------|------------------------|--------------|
| T1 | $B \bullet 1 = B$ | T1′ | B + 0 = B | Identity |
| T2 | $B \bullet 0 = 0$ | T2′ | B + 1 = 1 | Null Element |
| T3 | $B \bullet B = B$ | T3′ | B + B = B | Idempotency |
| T4 | | $\overline{\overline{B}} = B$ | | Involution |
| T5 | $B \bullet \overline{B} = 0$ | T5' | $B + \overline{B} = 1$ | Complements |

Dual: Exchange: • and + 0 and 1



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Chapter 2 <32>



T1: Identity Theorem

- B 1 = B
- $\mathbf{B} + \mathbf{0} = \mathbf{B}$



T1: Identity Theorem

- $B \cdot 1 = B$
- $\mathbf{B} + \mathbf{0} = \mathbf{B}$





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Chapter 2 <34>

Switching Algebra

- Simplification of digital logic → connecting wires with a on/off switch
- X = 0 (switch open)
- X = 1 (switch closed)





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Chapter 2 <35>

Series Switching Network: AND

• Switching circuit in series performs AND



• 1 is connected to 2 iff A AND B are 1



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Chapter 2 <36>

Parallel Switching Network: OR

• Switching circuit in parallel performs OR



• 1 is connected to 2 if A **OR** B is 1



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Chapter 2 <37>

T1: Identity Theorem

- $B \cdot 1 = B$
- $\mathbf{B} + \mathbf{0} = \mathbf{B}$







T2: Null Element Theorem

- $B \cdot 0 = 0$
- B + 1 = 1





T2: Null Element Theorem

- $B \cdot 0 = 0$
- B + 1 = 1







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Chapter 2 <40>

T3: Idempotency Theorem

- $\mathbf{B} \cdot \mathbf{B} = \mathbf{B}$
- B + B = B



T3: Idempotency Theorem

- $\mathbf{B} \cdot \mathbf{B} = \mathbf{B}$
- B + B = B







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Chapter 2 <42>



T4: Involution Theorem

• $\overline{B} = B$



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Chapter 2 <43>



T4: Involution Theorem

• $\overline{B} = B$





T5: Complements Theorem

- $\mathbf{B} \cdot \overline{\mathbf{B}} = \mathbf{0}$
- $B + \overline{B} = 1$



T5: Complements Theorem

- B $\overline{B} = 0$
- $B + \overline{B} = 1$







Recap: Basic Boolean Theorems

| | Theorem | | Dual | Name |
|----|------------------------------|-------------------------------|------------------------|--------------|
| T1 | $B \bullet 1 = B$ | T1′ | B + 0 = B | Identity |
| Т2 | $B \bullet 0 = 0$ | T2′ | B + 1 = 1 | Null Element |
| Т3 | $B \bullet B = B$ | T3′ | B + B = B | Idempotency |
| T4 | | $\overline{\overline{B}} = B$ | | Involution |
| T5 | $B \bullet \overline{B} = 0$ | T5' | $B + \overline{B} = 1$ | Complements |



Boolean Theorems of Several Vars

| Number | Theorem | Name |
|--------|--|----------------|
| Т6 | $B \bullet C = C \bullet B$ | Commutativity |
| Т7 | $(B \bullet C) \bullet D = B \bullet (C \bullet D)$ | Associativity |
| Т8 | $B \bullet (C + D) = (B \bullet C) + (B \bullet D)$ | Distributivity |
| Т9 | B● (B+C) = B | Covering |
| T10 | $(B \bullet C) + (B \bullet \overline{C}) = B$ | Combining |
| T11 | $B \bullet C + (\overline{B} \bullet D) + (C \bullet D) =$ $B \bullet C + \overline{B} \bullet D$ | Consensus |



Boolean Theorems of Several Vars

| Number | Theorem | Name |
|--------|--|----------------|
| Т6 | $B \bullet C = C \bullet B$ | Commutativity |
| Т7 | $(B \bullet C) \bullet D = B \bullet (C \bullet D)$ | Associativity |
| Т8 | $B \bullet (C + D) = (B \bullet C) + (B \bullet D)$ | Distributivity |
| Т9 | B● (B+C) = B | Covering |
| T10 | $(B \bullet C) + (B \bullet \overline{C}) = B$ | Combining |
| T11 | $B \bullet C + (\overline{B} \bullet D) + (C \bullet D) =$ $B \bullet C + \overline{B} \bullet D$ | Consensus |

How do we prove these are true?



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Chapter 2 <49>

How to Prove Boolean Relation

- Method 1: Perfect induction
- Method 2: Use other theorems and axioms to simplify the equation
 - Make one side of the equation look like the other



Proof by Perfect Induction

- Also called: proof by exhaustion
- Check every possible input value
- If two expressions produce the same value for every possible input combination, the expressions are equal



Example: Proof by Perfect Induction

| Number | Theorem | Name |
|--------|-----------------------------|---------------|
| Т6 | $B \bullet C = C \bullet B$ | Commutativity |

| В | С | ВС | СВ | |
|---|---|----|----|--|
| 0 | 0 | | | |
| 0 | 1 | | | |
| 1 | 0 | | | |
| 1 | 1 | | | |
| | | | | |



Example: Proof by Perfect Induction

| Number | Theorem | Name |
|--------|-----------------------------|---------------|
| Т6 | $B \bullet C = C \bullet B$ | Commutativity |

| В | С | ВС | СВ | |
|---|---|----|----|--|
| 0 | 0 | 0 | 0 | |
| 0 | 1 | 0 | 0 | |
| 1 | 0 | 0 | 0 | |
| 1 | 1 | 1 | 1 | |



Boolean Theorems of Several Vars

| Number | Theorem | Name |
|--------|--|----------------|
| Т6 | $B \bullet C = C \bullet B$ | Commutativity |
| Т7 | $(B \bullet C) \bullet D = B \bullet (C \bullet D)$ | Associativity |
| Т8 | $B \bullet (C + D) = (B \bullet C) + (B \bullet D)$ | Distributivity |
| Т9 | B● (B+C) = B | Covering |
| T10 | $(B \bullet C) + (B \bullet \overline{C}) = B$ | Combining |
| T11 | $B \bullet C + (\overline{B} \bullet D) + (C \bullet D) =$ $B \bullet C + \overline{B} \bullet D$ | Consensus |



T7: Associativity

| Number | Theorem | Name |
|--------|---|---------------|
| Т7 | $(B \bullet C) \bullet D = B \bullet (C \bullet D)$ | Associativity |



T8: Distributivity

| Number | Theorem | Name |
|--------|---|----------------|
| Т8 | $B \bullet (C + D) = (B \bullet C) + (B \bullet D)$ | Distributivity |







| Number | Theorem | Name |
|--------|--------------|----------|
| Т9 | B● (B+C) = B | Covering |

Prove true by:

- Method 1: Perfect induction
- Method 2: Using other theorems and axioms



| Number | Theorem | Name |
|--------|--------------|----------|
| Т9 | B● (B+C) = B | Covering |

Method 1: Perfect Induction

| В | С | (B+C) | B(B+C) | |
|---|---|-------|--------|--|
| 0 | 0 | | | |
| 0 | 1 | | | |
| 1 | 0 | | | |
| 1 | 1 | | | |
| | | | | |
| | | | | |
| | | l | | |



| Number | Theorem | Name |
|--------|--------------|----------|
| Т9 | B● (B+C) = B | Covering |

Method 1: Perfect Induction

| В | С | (B+C) | B(B+C) | |
|-------|---|-------|--------|--|
| 0 | 0 | 0 | 0 | |
| 0 | 1 | 1 | 0 | |
| 1 | 0 | 1 | 1 | |
| 1 | 1 | 1 | 1 | |



| Number | Theorem | Name |
|--------|--------------|----------|
| Т9 | B● (B+C) = B | Covering |

Method 2: Prove true using other axioms and theorems.



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Chapter 2 <61>

| Number | Theorem | Name |
|--------|--------------|----------|
| Т9 | B● (B+C) = B | Covering |

Method 2: Prove true using other axioms and theorems.

 $B \bullet (B+C) = B \bullet B + B \bullet C$ T8: Distributivity

 $= \mathbf{B} + \mathbf{B} \cdot \mathbf{C}$

 $= B \cdot (1 + C)$

 $= \mathbf{B} \cdot (\mathbf{1})$

 $\equiv \mathbf{B}$

- T3: Idempotency
- T8: Distributivity
- T2: Null element

| T1: | Identity |
|---------|----------|
| Chapter | 2 <62> |



T10: Combining

| Number | Theorem | Name |
|--------|--|-----------|
| T10 | $(B \bullet C) + (B \bullet \overline{C}) = B$ | Combining |

Prove true using other axioms and theorems:



T10: Combining

| Number | Theorem | Name |
|--------|--|-----------|
| T10 | $(B \bullet C) + (B \bullet \overline{C}) = B$ | Combining |

Prove true using other axioms and theorems: $B \cdot C + B \cdot \overline{C} = B \cdot (C + \overline{C})$ T8: Distributivity $= B \cdot (1)$ T5': Complements = B T1: Identity



T11: Consensus

| Number | Theorem | Name |
|--------|--|-----------|
| T11 | $(B \bullet C) + (\overline{B} \bullet D) + (C \bullet D) =$ $(B \bullet C) + \overline{B} \bullet D$ | Consensus |

Prove true using (1) perfect induction or (2) other axioms and theorems.



Recap: Boolean Thms of Several Vars

| Number | Theorem | Name |
|--------|--|----------------|
| Т6 | $B \bullet C = C \bullet B$ | Commutativity |
| Т7 | $(B \bullet C) \bullet D = B \bullet (C \bullet D)$ | Associativity |
| Т8 | $B \bullet (C + D) = (B \bullet C) + (B \bullet D)$ | Distributivity |
| Т9 | B● (B+C) = B | Covering |
| T10 | $(B \bullet C) + (B \bullet \overline{C}) = B$ | Combining |
| T11 | $B \bullet C + (\overline{B} \bullet D) + (C \bullet D) =$ $B \bullet C + \overline{B} \bullet D$ | Consensus |



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Chapter 2 <66>

Boolean Thms of Several Vars: Duals

| # | Theorem | Dual | Name |
|-----|--|--|----------------|
| T6 | $B \bullet C = C \bullet B$ | B+C = C+B | Commutativity |
| Т7 | $(B \bullet C) \bullet D = B \bullet (C \bullet D)$ | (B + C) + D = B + (C + D) | Associativity |
| T8 | $B \bullet (C + D) = (B \bullet C) + (B \bullet D)$ | $B + (C \bullet D) = (B+C) (B+D)$ | Distributivity |
| Т9 | B • (B+C) = B | $B + (B \bullet C) = B$ | Covering |
| T10 | $(B \bullet C) + (B \bullet \overline{C}) = B$ | $(B+C) \bullet (B+\overline{C}) = B$ | Combining |
| T11 | $(B \bullet C) + (B \bullet \overline{D}) + (C \bullet D) =$ $(B \bullet C) + (B \bullet \overline{D})$ | $(B+C) \bullet (B+\overline{D}) \bullet (C+D) =$ $(B+C) \bullet (B+\overline{D})$ | Consensus |
| | | | |

Dual: Replace: • with + 0 with 1



Boolean Thms of Several Vars: Duals

| Theorem | Dual | Name |
|--|---|--|
| $B \bullet C = C \bullet B$ | B+C = C+B | Commutativity |
| $(B \bullet C) \bullet D = B \bullet (C \bullet D)$ | (B + C) + D = B + (C + D) | Associativity |
| $B \bullet (C + D) = (B \bullet C) + (B \bullet D)$ | $B + (C \bullet D) = (B+C) (B+D)$ | Distributivity |
| B • (B+C) = B | $B + (B \bullet C) = B$ | Covering |
| $(B \bullet C) + (B \bullet \overline{C}) = B$ | $(B+C) \bullet (B+\overline{C}) = B$ | Combining |
| $(B \bullet C) + (B \bullet \overline{D}) + (C \bullet D) =$ $(B \bullet C) + (B \bullet \overline{D})$ | $(B+C) \bullet (B+\overline{D}) \bullet (C+D) =$ $(B+C) \bullet (B+\overline{D})$ | Consensus |
| | Theorem $B \cdot C = C \cdot B$ $(B \cdot C) \cdot D = B \cdot (C \cdot D)$ $B \cdot (C + D) = (B \cdot C) + (B \cdot D)$ $B \cdot (B + C) = B$ $(B \cdot C) + (B \cdot \overline{C}) = B$ $(B \cdot C) + (B \cdot \overline{D}) + (C \cdot D) =$ $(B \cdot C) + (B \cdot \overline{D})$ | TheoremDual $B \bullet C = C \bullet B$ $B + C = C + B$ $(B \bullet C) \bullet D = B \bullet (C \bullet D)$ $(B + C) + D = B + (C + D)$ $B \bullet (C + D) = (B \bullet C) + (B \bullet D)$ $B + (C \bullet D) = (B + C) (B + D)$ $B \bullet (B + C) = B$ $B + (B \bullet C) = B$ $(B \bullet C) + (B \bullet \overline{C}) = B$ $(B + C) \bullet (B + \overline{C}) = B$ $(B \bullet C) + (B \bullet \overline{D}) + (C \bullet D) = (B + C) (B + \overline{D}) \bullet (C + D) = (B + C) \bullet (B + \overline{D})$ |

Dual: Replace: • with + 0 with 1

Warning: T8' differs from traditional algebra: OR (+) distributes over AND (•)



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Chapter 2 <68>

Boolean Thms of Several Vars: Duals

| # | Theorem | Dual | Name |
|-----|--|--|----------------|
| T6 | $B \bullet C = C \bullet B$ | B+C = C+B | Commutativity |
| T7 | $(B \bullet C) \bullet D = B \bullet (C \bullet D)$ | (B + C) + D = B + (C + D) | Associativity |
| T8 | $B \bullet (C + D) = (B \bullet C) + (B \bullet D)$ | $B + (C \bullet D) = (B+C) (B+D)$ | Distributivity |
| Т9 | B • (B+C) = B | $B + (B \bullet C) = B$ | Covering |
| T10 | $(B \bullet C) + (B \bullet \overline{C}) = B$ | $(B+C) \bullet (B+\overline{C}) = B$ | Combining |
| T11 | $(B \bullet C) + (B \bullet \overline{D}) + (C \bullet D) =$ $(B \bullet C) + (B \bullet \overline{D})$ | $(B+C) \bullet (B+\overline{D}) \bullet (C+D) =$ $(B+C) \bullet (B+\overline{D})$ | Consensus |

Axioms and theorems are useful for *simplifying* equations.

